Zagreb, Cloud analysis

Container technologies

Davor Popović · Marko Ratkaj





About us

• Delivering solutions based on Linux, OpenWrt and Yocto

- Focused on software in network edge and CPEs
- Continuous participation in Open Source projects
- www.sartura.hr



Introduction to GNU/Linux







• First Linux prototypes publicly released in 1991

- https://github.com/zavg/linux-0.01
- Version 1.0 release on March 14, 1994
- Written from the ground up



- Linux = operating system kernel
- GNU/Linux distribution = kernel + userspace (Ubuntu, Arch Linux, Gentoo, Debian, OpenWrt, *Mint*, ...)
- Userspace = set of libraries + system software
 - Libraries = resources for other software
 - System software = platform/environment for other software to run (operating system)







- Kernel = all are based on the Linux kernel
- Folder hierarchy = all follow similar folder hierarchy (minor differences between distros)
- Init systems = all have init systems (most use systemd, some use SystemV (sysV) or other (openRC)
- Package managers = all have package managers, but they differ between distributions (*apt, pacman, emerge, yum*, etc.)







(GNU/Linux) kernel

- Core part of an operating system which manages system resources
 - CPU
 - RAM
 - Input/Output devices (keyboards, mouse, disk drives, printers, USB devices, network adapters, and display devices)
- The first program loaded into RAM by the bootloader (BIOS, U-Boot...)
- Acts as a layer between software/applications (userspace) and hardware



Linux kernel

- Operating systems have two spaces of operation:
 - · Kernel space protected memory space and full access to the device's hardware
 - Userspace space in which all other application run
 - Has limited access to hardware resources
 - Accesses hardware resources via kernel
 - Userspace applications invoke kernel services with system calls



Example: opening a file

- Application wants to open a file → invokes library call for open → invokes system call for open → kernel opens a file from disk signaling the success result back to application in the same ladder
- A small programming detail for C language:
 - For opening files fopen(const char *path, const char *mode) from stdio.h is used(glibc)
 - Programs invoke this call from the source code with necessary parameters
 - fopen contains a system call for opening a file which is a function from open family of the system calls





Example: opening a file

• Example taken from https://github.com/freebsd/freebsd/blob/master/lib/libc/stdio/fopen.c

```
1 FILE *
2 fopen(const char * __restrict file, const char * __restrict mode)
3 {
4 ...
5 if ((f = _open(file, oflags, DEFFILEMODE)) < 0) {
6 fp->_flags = 0;^^1^^1^1/* release */
7 return (NULL);
8 }
9 ...
```

• system call open returns a pointer to structure FILE back to the caller



• Every other interaction follows the same principle but with a different system call



Linux kernel

- System calls provide the following services from the kernel to applications:
 - 1. Process creation and management
 - 2. Main memory management
 - 3. File Access, directory and file system management
 - 4. Device handling (I/O)
 - 5. Protection
 - 6. Networking, etc.



Linux kernel

- System calls can be grouped in the following types:
 - 1. Process control: end, abort, create, terminate, allocate and free memory.
 - 2. File management: create, open, close, delete, read file etc.
 - 3. Device management
 - 4. Information maintenance
 - 5. Communication





	User applications	E.g. bash, LibreOffice, G	IMP, Blender, Mozilla Fire	fox, etc.			
User mode	Low-level system components	System daemons: systemd, runit, logind, networkd, PulseAudio, 	Windowing system: X11, Wayland, SurfaceFlinger (Android)	Other libraries: GTK+ FLTK, GNUstep, etc.	, Qt, EFL, SDL, SFML,	Graphics: Mesa, AMD Catalyst,	
	C standard library	Up to 2000 subroutines fopen(), calloc(),	depending on C library ()	glibc, musl, uClibc, bionic) (a	open(), exec(), sbrk(), socket(),	
		About 380 system calls	(stat, splice, dup, :	read, open, ioctl, write	e, mmap, close, exit,	etc.)	
Kernel mode	Linux Kernel	Process scheduling subsystem	IPC subsystem	Memory management subsystem	Virtual files subsystem	Network subsystem	
		Other components: ALS SELinux, TOMOYO, App/	A, DRI, evdev, LVM, device Armor, Smack	e mapper, Linux Network Sch	eduler, Netfilter Linux Sec	urity Modules:	
		Hardware (CPU, main memory, data	a storage devices, etc.)			
		TAB	LE 1 Layers w	ithin Linux/			
_							
5							

Virtualization





- Virtualization = creating a virtual version of hardware/software
- Term coined in 1960s with mainframe computers
- What can be virtualized? basically everything
 - Hardware (with emulation)
 - Storage devices (with pooling of physical storage from multiple storage devices into a single logical storage device)
 - Can be applied on HDD, RAM
 - Networking, processes (using various kernel features that allow resource separation)
 - These are operating system features rather than applications
 - Software/applications
 - Running applications in separate machines/containers and making them portable
 - Running different services in separate machines/containers (cloud-based applications)





Virtualization Concepts

Two virtualization concepts:

- Hardware virtualization (full/para virtualization)
 - Emulation of complete hardware (virtual machines VMs)
 - VirtualBox, QEMU, etc.
- Operating system level virtualization
 - Utilizing kernel features for running more than one userspace instance
 - Each instance is isolated from the rest of the system and other instances 1000
 - Method for running isolated processes is called a *container*
 - Docker, LXC, Solaris Containers, Microsoft Containers, rkt, etc.





Virtual machines use hypervisors (virtual machine managers – VMM)

- Allows multiple guest operating systems (OS) to run on a single host system at the same time
- Responsible for resource allocation each VM uses the real hardware of the host machine but presents system components (e.g. CPU, memory, HDD, etc.) as their own
- Two types of hypervisors (VMMs)
 - Type 1
 - Native or bare-metal hypervisors running directly on host's hardware and controlling the resources for Vms (Microsoft Hyper-V, VMware ESXi, Citrix XenServer)
 - Type 2
 - Hosted hypervisors running within a formal operating system environment.
 - Host OS acts as a layer between hypervisor and hardware.







- Containers do not use hypervisors
- Containers sometimes come with *container managers*
 - Used for managing containers rather than resource allocation

• Containers use direct system calls to the kernel to perform actions

• Kernel is shared with the host





- Idea behind containers is to pack the applications with all their dependencies and run them in an environment that is isolated from the host
- Two types of containers:
 - Full OS containers contain full root file system of the operating system
 - · Meant to run multiple applications at once
 - Provide full userspace isolation
 - LXC, BSD jails, OpenVZ, Linux VServer, Solaris Zones
 - Application containers contain an application which is isolated from the rest of the system (*sandboxing*)
 - Application behaves at runtime like it is directly interfacing with the original operating system and all the resources managed by it
 - Docker, rkt



Parameter	VMs	Containers
Size	Few GBs	Few MBs
Structure	Full contained environment	Rely on underlying OS
Resources	Contains full OS with no dependencies on the underlying C (e.g. Windows running on Linux and vice-versa)	OS Rely on underlying OS
Boot time	Few second overhead	Millisecond overhead
	TABLE 2 VMs vs containers - Differences	
5		

Libraries			
Applications			
Binaries	Applications rootfs		
Guest OS Guest OS Guest	Container Container Conta	iner	
Hypervisor			
Host OS	Host OS		
Hardware	Hardware		
Virtualization	Containers		
FIGURE 3 Virtualiza	ation vs Containers		

Why use virtualization?

- Cost effective, resource savings
 - Multiple machines can be virtualized on a single machine
- Management
 - Everything can be managed from a single point, usually using a management software for virtual machines and/or containers
- Maintenance
 - Once deployed, machines can be easily switched for new machines if needed in the future





Linux Containers (LXC)





- LXC is a userspace interface for the GNU/Linux kernel containment features
 - Allows operating system level virtualization on GNU/Linux systems
- In-between *chroot* and complete VM
 - Sometimes referred to as chroot-on-steroids
- Does not depend on hardware support for virtualization
 - · Ideal for containerization/virtualization on embedded devices
- Configurable as a full feature file system (rootfs) or minimized for running single applications
- Relies *heavily* on kernel features



Kernel features

• LXC uses the following kernel features:

- Namespaces
 - ipc, uts, mount, pid, network and user
- Apparmor and SELinux profiles
- Seccomp policies
- Kernel capabilities
- CGroups
- Chroots (using pivot_root)





Namespaces

- Lightweight process virtualization
- Namespaces = way of grouping items under the same designation
 - Kernel feature which organizes resources for processes
 - One process sees one set of resources
 - Another process sees another set of resources
 - One process cannot see other processes' set of resources
 - Each process has its own namespace a set of resources uniquely allocated for that process
 - Namespaces allow processes to see the same parts of the system differently



Namespaces

• GNU/Linux kernel supports the following types of namespaces:

- network
- uts
- PID
- mount
- user

S



Namespaces - Network

- Logical copy of the network stack containing its own network features (routes, firewall rules and network devices)
- Initial network namespace:
 - Initial network namespace = loopback device + physical devices + networking tables + ...
- Every other newly created namespace contains only a loopback device
 - New network namespace instance = loopback device
- Every other device must be added manually
- To have network devices every container must be configured via configuration files





Namespaces - Network

• Network namespaces are completely isolated from each other

- · Host network namespace and container namespaces do not see each other
- To get to the network, network namespaces need to be connected
- E.g. the inside interface in the container must be connected to a physical interface on the host to go online









Namespaces - Network

- Connecting network namespaces
 - Done with a virtual interface called <code>veth</code>
 - veth is actually a "virtual cable" handled by the kernel one end is "plugged" into a container, the other end into the host interface
 - Done with a virtual interface called macvlan
 - macvlan is a virtual interface which has its own MAC address and can be directly associated with a physical interface of the host







FIGURE 4 Network namespaces - veth



Namespaces – PID

- Process = current program in execution (set of instructions that are currently occupying CPU cycles)
- PID = way of organizing processes inside of the kernel by a unique number each process is assigned a number
 - Special case regarding process with PID 1 the first process started when OS is booted up (called init system) which starts all other processes
 - GNU/Linux supports multiple init systems *systemd*, *systemV*, *openRC*, etc.
 - Processes communicate with signals (e.g. kill signal)
- PID namespaces kernel feature enables isolating PID namespaces
 - Processes in different PID namespaces can have the same PID





Namespaces – PID

- Kernel creates namespace one NS1
- Kernel creates namespace two NS2
- NS1 contains PID 1 and all other PIDs
- NS2 contains PID1 and all other PIDs
- NS1 can see PID 1 of NS2 but with some other PID (e.g. PID 10001) and all other PIDs of NS2
- NS2 can not see PIDs from NS1
- This way, isolation is achieved from these namespaces processes inside NS2 are only functional if NS2 is isolated from NS1
 - In NS2 a process cannot send signal (e.g. kill) to a host machine



CGroups

- PID namespace allows processes to be grouped together in isolated environments
- Group of processes (or a single process) needs access to certain hardware components
 - E.g. RAM, CPU, ...
- Kernel provides the control groups (CGroups) feature for limiting how process groups access and use these resources







CGroups

- 4 main purposes
 - Limiting resources = groups can be set to not exceed a pre-configured memory limit (i.e. this group of processes can access X MB of RAM)
 - Prioritization = some groups may get a larger share of CPU utilization (i.e. this group of processes can utilize 43% if CPU 1, while another group only 5%)
 - Accounting = measures a group's resource usage (i.e. this group of processes has been using 5% of CPU)
 - Useful for statistics
 - Control = used for freezing, snapshoting/checkpointing and restarting





CGroups

- CGroups uses a Virtual File System
 - /sys/cgroups
 - All CGroups actions are performed via file system actions





- Install LXC package with your distribution's package manager
- Kernel must have all the required features enabled
 - Some distributions may come with a complete required feature set (e.g. Ubuntu)
 - · Some distributions do not come with a complete required feature set
 - Some distributions require manual kernel configurations (e.g. Gentoo)
 - After installing LXC, it is required to check that the kernel has everything it needs with lxc-checkconfig command





- Each container needs its own configuration file
- Each container needs its own root file system which will run in side of the container
 - Root file system contains all the necessary libraries, applications and environment settings
 - Root file system needs to be either manually prepared or downloaded from remote online repositories
- Place configuration file and root file system on the same location:
 - /var/lib/lxc/<container_name>/





Configuring the container

- Container configuration defines the following components:
 - · Capabilities what the container is allowed to do from an administrative perspective
 - Cgroups which resources of the host are allowed for the container (e.g. configuring which devices can the container use)
 - Mount namespaces which of the host folders/virtual file systems will be allowed for mounting inside the container (virtual file systems such as proc or sys)
 - Network namespaces which devices will be created inside the container and how they connect to the outside network



- First part of the file handles capabilities
- A list of all the capabilites dropped (not allowed) for the container
- Usually best to consult with man pages http://man7.org/linux/manpages/man7/capabilities.7.html

9

lxc.cap.drop	=	mac_admin		
lxc.cap.drop	=	mac_override		
lxc.cap.drop	=	sys_admin		
lxc.cap.drop	=	sys_boot		
lxc.cap.drop	=	sys_module		
lxc.cap.drop	-	sys_nice		
lxc.cap.drop	=	sys_pacct		
lxc.cap.drop	=	sys_ptrace		
lxc.cap.drop	=	sys_rāwio		
lxc.cap.drop	=	sys_resource		
lxc.cap.drop	=	sys_time		
lxc.cap.drop	=	sys_tty_config		
lxc.cap.drop	=	syslog		
lxc.cap.drop	=	wake_alarm		







- It uses traditional Linux desigantions for devices (try running ls -1 /dev which will list all the devices with corresponding major:minor numbers)
- E.g. bold entry is for console on PC

<pre>lxc.cgroup.devices.deny = a</pre>
<pre>lxc.cgroup.devices.allow = c 1:1 rwm</pre>
<pre>lxc.cgroup.devices.allow = c 1:3 rwm</pre>
<pre>lxc.cgroup.devices.allow = c 1:5 rwm</pre>
lxc.cgroup.devices.allow = c 5:1 rwm
<pre>lxc.cgroup.devices.allow = c 5:0 rwm</pre>
<pre>lxc.cgroup.devices.allow = c 4:0 rwm</pre>
<pre>lxc.cgroup.devices.allow = c 4:1 rwm</pre>
<pre>lxc.cgroup.devices.allow = c 1:9 rwm</pre>
<pre>lxc.cgroup.devices.allow = c 1:8 rwm</pre>
<pre>lxc.cgroup.devices.allow = c 1:11 rwm</pre>
<pre>lxc.cgroup.devices.allow = c 136:* rwm</pre>
<pre>lxc.cgroup.devices.allow = c 5:2 rwm</pre>
<pre>lxc.cgroup.devices.allow = c 254:0 rwm</pre>
<pre>lxc.cgroup.devices.allow = c 10:200 rwm</pre>



Some metadata about the container
Where is the rootfs located
Hostname of the container
What to mount from the host
/proc and /sys

8

<pre># Distribution c lxc.arch = x86_6 # Container spec lxc.rootfs.path lxc.uts.name = o # Mount entries lxc.mount.entry nouid 0 0 lxc.mount.entry</pre>	onfiguration 4 ific configuration dir:/var/lib/lxc/openwrt/rootfs penwrt = /prod prod /prod nodev,hoexec,



- Network namespace configuration
- We can read this as follows:
 - Create eth0 device inside a container
 - Use a veth (virtual cable) to connect this eth0 from container to lxcbr0 interface (a bridge interface) on the host
- It can be configured in many different ways depends on the use case

5

- Once configuration and root file system are ready issue lxc-ls
 - If the container is configured properly and its root file system is valid, the container should appear on the list
- Start the container with

lxc-start -n <container_name>

- There will be no output, but the container should start
- Check that the container is running

lxc-info -n <container_name>

 Container is running in the background, and we can now run applications inside of the container



- Entering the shell of the container (attaching inside of the container)
 - lxc-attach -n <container_name>
- From the shell, it is possible to do everything as on host GNU/Linux system
- To exit, type exit
- To stop the container
- lxc-stop -n <container_name>
- This demonstration is a simple case of a single container created by root user and with no particular functionalities so what can be done with the container?









- Containers generally have two purposes
 - Development when developing, containers allow developers to share the same environment once a container is set, it can be deployed on many machines with exactly the same libraries and user space
 - 2. Deployment containers allow quick deployment. It is possible to run applications in the cloud inside of the container allowing external access to them.
 - E.g. running a container with web servers allowing access from the outside of the network to each web server. If the security is compromised, in theory the attacker would be inside of the container isolated from the rest of the cloud. This method allows running multiple web servers on a single physical machine all listening on the same port. The additional effort is in creating firewall rules between the outside network and containers.







• Cloud machines allow running tens of thousands containers at once

- All hell breaks loose when all these containers need to be administered and supported in the future
- LXD was created to answer challenges with managing huge numbers of containers









LXD

- LXD = container manager, useful when running and configuring huge amounts of Linux containers
- Concept
 - Server + client side (communicating over REST API)
 - Accessible locally and remotely over network
 - Command line tool for working with containers
- Supports the full LXC feature set
 - By default, LXD creates unprivileged containers (what we demonstrated is the creation of privileged containers by the root user which might have some security issues)



- Scalable allows creating containers over thousands of hardware nodes (PC, servers, embedded)
- Intuitive simple API and CLI
- Image-based huge variety of distributions available over networkd
- Allows live migrations
- Resources control (CGroups CPU, memory, I/O, ...)
- Device pass-through (allowing USB, GPU, character devices, ...)
- Networking
- Storage (storage pools and volumes)





LXD - Prerequisites

- Initialized LXD daemon
- Root file system and metadata
 - · Metadata = data about the container
- Container image
 - · Image from which the container will be created
 - Image = rootfs + metadata
- Container profile
 - Basic container configuration



LXD init

• Configuring the LXD daemon

lxd init

S



Prepare rootfs

• Create gentoo directory inside the home directory

cd ~

mkdir gentoo

• Copy compressed root file system on that location

```
cp gentoo-rootfs.tar.gz ~/gentoo
```

- In the same directory, create metadata file for the container
 - · Metadata describes basic information about the container
 - Can be written in YAML format or JSON (examples below use YAML)





Import rootfs and metadata as image

- Compress both image and metadata
 - tar cf gentoo-matadata.tar metadata.yaml
- Import compressed root file system and metadata into LXD
 - lxc image import gentoo-metadata.tar.gz gentoo-rootfs.tar.gz --alias GentooImage
- If everything went well, the image should appear on LXD image list
 - lxc image list











LXD

- At this point, the container is still not created only the image for the container is prepared (root file system + metadata)
- Containers do have default profiles but in practice each container needs a profile with different configuration
- Profile = initial configuration for the container (networking, storage, hostname, ...)
- The example below creates a simple profile with one network interface connected to the host physical interface with macvlan virtual interface (see Figure 5)





Prepare container profile

• Create minimal YAML file to define the container profile: vim gentoo-profile.yaml



LXD

- At this point the profile is not attached to any container, and is actually just a file
- First step is to create a profile for LXD and apply the YAML file with the profile
- 2
- 1xd profile create Gentoo-profile
- lxd profile edit Gentoo-profile < gentoo-profile.yaml</pre>
- This profile can be used over n number of containers









LXD

- Next step is to apply the profile to a container
- First, container must be created from the image

lxc init GentooImage GentooContainer

• Then, apply the profile to the initialized container

lxc profile apply Gentoo-Profile GentooContainer



Verifying the process

- To verify what has been done (and if it went OK)
 - Checking images
 - lxc image list
 - Checking containers

lxc ls

Checking available profiles

lxc profile list

S

Verifying the process

• If necessary, profiles can be modified on the fly and all changes applied in real time

Checking a specific profile

lxc profile show Gentoo-profile

Modifying a specific profile

lxc profile edit Gentoo-profile



Starting the container

• The container is ready to start at this point

lxc start GentooContainer

- How does this all fit together?
 - Inspect htop
 - Network namespace



Starting the container - next steps

• To attach inside the container, execute bash lxc exec GentooContainer /bin/bash
lxc exec GentooContainer /bin/bash
From this shell we can do everything as regular Linux users
Any other application can be run in the same way
lxc exec GentooContainer /bin/bash
• With this principle different servers and applications can be run inside the container to isolate them from the rest of the host system
Once a container is stopped (or destroyed), applications have no power

SDN – practical application of Containers

- The core network component = virtual router
- Virtual router = LXC container running Linux with networking set for router performance
 - What is a router?
 - · Forwards packets between two different networks
 - One side is LAN one side is WAN with NAT in-between







• Each router needs to have profiles defined to support two sides of the network (LAN and WAN)

• How can this be done?

L	eth0:
2	name: eth-wan
3	nictype: macvlan
1	parent: eth2
5	type: nic
5	eth1:
7	name: eth-lan
3	nictype: macvlan
9	parent: tapl
)	type: nic













- How to configure network inside the container? Remember systemd?
- LAN should behave as a DHCP server
- WAN should behave as a DHCP client
- Network Addres Translation (NAT) in between
 - · Keeps track who is going where on the Internet
 - Achieved with Linux firewall (iptables)



Container technologies

davor.popovic@sartura.hr · marko.ratkaj@sartura.hr



info@sartura.hr · www.sartura.hr

